- 1. Suppose that  $\varphi$  and  $\psi$  are formulas in propositional logic,  $\varphi$  is not a contradiction,  $\psi$  is not a tautology, and  $(\varphi \to \psi)$  is a tautology. Show that there is a propositional formula  $\theta$  such that every propositional symbol that appears in  $\theta$  also appears in both  $\varphi$  and  $\psi$  and such that  $(\varphi \to \theta)$  and  $(\theta \to \psi)$  are tautologies.
- 2. Show that  $(\mathbb{Z}, +)$  is minimal, but not prime.
- 3. Assume that  $\mathfrak{M}$  and  $\mathfrak{N}$  are elementarily equivalent structures for a countable language. Assume that  $\mathfrak{M}$  is countable. Show that  $\mathfrak{M}$  embeds into any ultrapower of  $\mathfrak{N}$  by a nonprincipal ultrafilter.
- 4. *I* is an index set if and only if  $e \in I$  and  $W_e = W_n$  implies  $n \in I$ , where  $(W_i : i \in \mathbb{N})$  is the canonical indexing of recursively enumerable sets. Let *I* be a recursively enumerable index set, and let  $C = \{W_e : e \in I\}$ . Show:
  - (a) If  $W_e \in C$  then there is a finite subset F of  $W_e$  such that  $F \in C$ .
  - (b) If  $F \in C$  and  $F \subseteq W_n$  then  $W_n \in C$ .
  - (c) There is a recursively enumerable set S of canonical finite sets such that for all  $e, e \in I$  if and only if for some  $F \in S, F \subseteq W_e$ .
- 5. Show that the index set  $\{e: W_e \text{ is infinite}\}\$  is not recursive relative to  $\emptyset'$ .
- 6. Let  $\varphi$  be a  $\Pi_1^0$  sentence st  $PA \vdash (\varphi \rightarrow Con(PA))$ . Show there is a sentence  $\psi$  such that  $PA \vdash (\varphi \leftrightarrow Con(PA \cup \{\psi\}))$ .

Hint: for  $\varphi = \forall w \theta(w)$ , where  $\theta$  is limited, show there is a  $\Sigma_1^0$  sentence  $\psi$  such that  $PA \vdash (\psi \leftrightarrow (\exists p)(\forall w < p)[\theta(w) \land "p \text{ is a proof of } \neg \psi \text{ from } PA"])$ 

- 7. Let M be a structure and  $a \in M$  an element. Assume that there are infinitely many elements in M which have the same type as a (over  $\emptyset$ ). Show that there is an elementary extension  $M \prec N$  and an automorphism  $\sigma$  of N such that  $a, \sigma(a), \sigma^2(a), \sigma^3(a), \ldots$  are pairwise distinct.
- 8. Let L be a finite relational language and M an L-structure. Assume that M admits elimination of quantifiers in L. Show that the theory of M is  $\omega$ -categorical.